

$$C = \log_2 \left[1 + \exp\left(\frac{\alpha - h}{1 - \alpha}\right) \right]$$

Example 2.23

If π_1 and π_2 are channel matrices of discrete memoryless channels k_1 and k_2 respectively, the sum of k_1 and k_2 is defined as the channel whose matrix is

$$\begin{bmatrix} \pi_1 & 0 \\ 0 & \pi_2 \end{bmatrix}$$

Thus the sum channel may be operated by choosing one of the individual channels and transmitting a digit through it. The input or output symbol always identifies the particular channel used. If C_i is the capacity of K_i , $i = 1, 2$ and C is the capacity of the sum show that

$$2C = 2^{C_1} + 2^{C_2}.$$

Solution

Let the input alphabet of k_1 and k_2 be x_1, \dots, x_r and x_{r+1}, \dots, x_M respectively; let the output alphabets of k_1 and k_2 be y_1, \dots, y_s and $y_{s+1}, \dots,$

y_N . If $p(x)$ is any input distribution, and $P = \sum_{i=1}^r p(x_i)$, then

$$\begin{aligned} H(X) &= - \sum_{i=1}^M p(x_i) \log p(x_i) \\ &= -p \log p - (1-p) \log(1-p) + pH_1(X) + (1-p) H_2(X) \end{aligned}$$

(By grouping axiom)

where
$$H_1(X) = - \sum_{i=1}^r \frac{p(x_i)}{p} \log \frac{p(x_i)}{p}$$

is the input uncertainty of k_1 under the distribution

$$\{p(x_i)/p, i = 1, \dots, r\}$$

and
$$H_2(X) = - \sum_{i=r+1}^M \frac{p(x_i)}{1-p} \log \frac{p(x_i)}{1-p}$$

is the input uncertainty of k_2 under the distribution

$$\{p(x_i)/(1-p), i = r+1, \dots, M\}.$$

[Alternately,
$$H(X) = - \sum_{i=1}^r p \frac{p(x_i)}{p} \log \left[\left(\frac{p(x_i)}{p} \right) p \right]$$

$$- \sum_{i=r+1}^M \frac{(1-p)p(x_i)}{(1-p)} \log \left[\frac{p(x_i)}{(1-p)} (1-p) \right]$$

leading to the same expression as above.]

Now,
$$H(X/Y) = - \sum_{i=1}^M \sum_{j=1}^N p(x_i, y_j) \log p(x_i / y_j)$$

$$= - \sum_{i=1}^r \sum_{j=1}^s p \frac{p(x_i)}{p} p(y_j / x_i) \log p(x_i / y_j)$$

$$- \sum_{i=r+1}^M \sum_{j=s+1}^N (1-p) \left[\frac{p(x_i)}{1-p} \right] p(y_j / x_i) \log p(x_i / y_j)$$

Now if $1 \leq i \leq r$, $1 \leq j \leq s$, the conditional probability.

$$p(x_i / y_j) = \frac{p(x_i)p(y_j / x_i)}{p(y_j)} = \frac{p[p(x_i)/p]p(y_j / x_i)}{p \sum_{i=1}^M \frac{p(x_i)}{p} \cdot p(y_j / x_i)}$$

$$= p\{X_1 = x_i / Y_1 = y_j\}$$

that is, the conditional probability that the input of k_1 is x_i given that the output of k_1 is y_j , under the distribution $\{p(x_i)/p, i = 1, 2, \dots, r\}$.

Thus

$$H(X/Y) = pH_1(X/Y) + (1-p)H_2(X/Y)$$

and

$$I(X; Y) = H(X) - H(X/Y).$$

Therefore, $H(X/Y) = H(p, 1-p) + p I_1(X; Y) + (1-p) I_2(X; Y)$

where the subscript i denotes that the indicated quantity is calculated for k_i under the appropriate input distribution. For a given p , we are completely free to choose the probabilities $p(x_i)/p$, $i = 1, \dots, r$ and $p(x_i)/(1-p)$, $i = r+1, \dots, M$.

We can do no better than to choose the probabilities so that

$$I_1(X; Y) = C_1, \quad I_2(X; Y) = C_2.$$

Thus it remains to maximize

$$H(p, 1-p) + pC_1 + (1-p)C_2,$$

Differentiating we obtain,

$$-1 - \log p + 1 + \log(1-p) + C_1 - C_2 = 0$$

$$p = \frac{2^{C_1}}{2^{C_1} + 2^{C_2}}$$

$$\begin{aligned} \text{Hence, } C &= H\left(\frac{2^{C_1}}{2^{C_1} + 2^{C_2}}, \frac{2^{C_2}}{2^{C_1} + 2^{C_2}}\right) + \frac{C_1 2^{C_1} + C_2 2^{C_2}}{2^{C_1} + 2^{C_2}} \\ &= \frac{2^{C_1}}{2^{C_1} + 2^{C_2}} \log(2^{C_1} + 2^{C_2}) + \frac{2^{C_2}}{2^{C_1} + 2^{C_2}} \log(2^{C_1} + 2^{C_2}) \end{aligned}$$

Therefore $C = \log(2^{C_1} + 2^{C_2})$

Example 2.24

Evaluate the capacity of the channel whose matrix is given as

$$\begin{bmatrix} 1-\beta & \beta & 0 \\ \beta & 1-\beta & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Solution

The given channel matrix is

$$\left[\begin{array}{cc|c} 1-\beta & \beta & 0 \\ \beta & 1-\beta & 0 \\ \hline 0 & 0 & 1 \end{array} \right]$$

where

$$\pi_1 = \begin{bmatrix} 1-\beta & \beta \\ \beta & 1-\beta \end{bmatrix}$$

$$\pi_2 = [1],$$

referring to the previous problem.

We have,

$$2^C = 2^{C_1} + 2^{C_2}$$

Here,

$$C_1 = \log m - h$$

$$= \log_2 2 - H(\beta, 1-\beta)$$

$$C_1 = 1 - H(\beta, 1-\beta)$$

$$C_2 = 0$$

Therefore,

$$2^C = 2^{1-H(\beta, 1-\beta)} + 2^0$$

$$C = \log [1 + 2^{1-H(\beta, 1-\beta)}]$$

Example 2.25

Evaluate the channel capacity of the channel whose matrix is given to be

$$\begin{bmatrix} \frac{1-p}{2} & \frac{1-p}{2} & \frac{p}{2} & \frac{p}{2} \\ \frac{p}{2} & \frac{p}{2} & \frac{1-p}{2} & \frac{1-p}{2} \end{bmatrix}$$

Solution

This is a symmetric channel.

$$\text{Channel capacity } C = \log_2 4 + p \log \frac{1}{2} p + (1-p) \log \frac{1}{2} (1-p)$$

or

$$C = 1 - H(p, 1-p)$$

= capacity of a binary symmetric channel with error probability p .

Example 2.26

Evaluate the channel capacity of the channel whose matrix is

0 erase 1

$$\begin{matrix} 0 & \begin{bmatrix} 1-p & p & 0 \end{bmatrix} \\ 1 & \begin{bmatrix} 0 & p & 1-p \end{bmatrix} \end{matrix}$$

Solution

This is a Binary erasure channel (BEC).

If $p(X=0) = \alpha$,

then $H(X) = H(\alpha, 1-\alpha)$

$$H(X/Y) = p(\text{erase}) H(X/Y = \text{erase})$$

$$= p \cdot H(\alpha, 1-\alpha)$$

Thus $I(X; Y) = (1-p) H(\alpha, 1-\alpha)$,

which is a maximum for $\alpha = 1/2$.

Therefore

$$C = 1 - p.$$

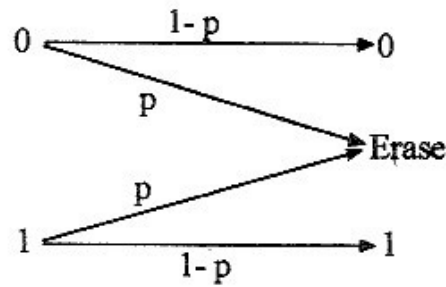


Fig. 2.16 Illustration of channels for problem 2.26

Example 2.27

A signal is bandlimited to 5 KHz, sampled at a rate of 10,000 samples per second, and quantized to 4 levels such that the sampled value can take values 0, 1, 2, 3 with probability p , p , $(1/2) - p$, $(1/2) - p$ respectively and the channel is characterized by the transition probability matrix

$$P = \{P_{ij}\} = P(X = x_i / Y = y_j)$$

$$Y = j$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 1/2 & 1/2 \end{bmatrix}$$

Find the capacity of the channel.

Solution

The channel model is shown below.

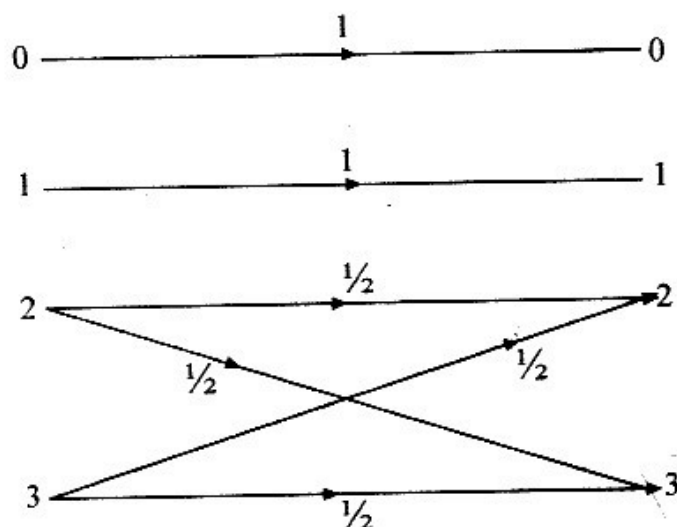


Fig. 2.17 Illustration of channels for problem 2.27

$$P(X = 0) = p$$

$$P(X = 1) = p$$

$$P(X = 2) = \frac{1}{2} - p$$

$$P(X = 3) = \frac{1}{2} - p$$

$$H(X) = - \sum_{i=0}^3 P(X = x_i) \log_2 [p(X = x_i)]$$

$$= -2p \log_2(p) + 2 \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} - p \right)$$

$$\begin{aligned} P(Y = 0) &= P(Y = 0 / X = 0) P(X = 0) + P(Y = 0 / X = 1) P(X = 1) \\ &\quad + P(Y = 0 / X = 2) P(X = 2) + P(Y = 0 / X = 3) P(X = 3) \\ &= 1(p) + 0 \left(\frac{1}{2} - p \right) + 0 \left(\frac{1}{2} - p \right) + 0(p) = p \end{aligned}$$

Similarly,

$$P(Y = 1) = p$$

$$P(Y = 2) = P(Y = 2 / X = 0) P(X = 0) + P(Y = 2 / X = 1) P(X = 1)$$

$$\begin{aligned}
 &+ P(Y = 2 / X = 2) P(X = 2) + P(Y = 2 / X = 3) P(X = 3) \\
 &= 0 \cdot p + 0 \left(\frac{1}{2} - p \right) + \frac{1}{2} \left(\frac{1}{2} - p \right) + \left(\frac{1}{2} - p \right) \frac{1}{2} \\
 &= \frac{1}{2} - p
 \end{aligned}$$

Similarly,

$$P(Y = 3) = \left(\frac{1}{2} - p \right)$$

$$\begin{aligned}
 P(X = 0, Y = 0) &= P(X = 0 / Y = 0) P(Y = 0) \\
 &= 1(p) = p
 \end{aligned}$$

$$\begin{aligned}
 P(X = 0, Y = 1) &= P(X = 0 / Y = 1) P(Y = 1) \\
 &= 0(p) = 0
 \end{aligned}$$

$$\begin{aligned}
 P(X = 0, Y = 2) &= P(X = 0 / Y = 0) P(Y = 2) \\
 &= 0 \left(\frac{1}{2} - p \right) = 0
 \end{aligned}$$

$$P(X = 0, Y = 3) = 0$$

$$P(X = 1, Y = 0) = 0$$

$$P(X = 1, Y = 1) = P(X = 1 / Y = 1) P(Y = 1) = 1 \cdot p = p$$

$$P(X = 1, Y = 2) = 0$$

$$P(X = 1, Y = 3) = 0$$

$$P(X = 2, Y = 0) = 0$$

$$P(X = 2, Y = 1) = 0$$

$$P(X = 2, Y = 2) = 0 = \frac{1}{2} \left(\frac{1}{2} - p \right)$$

$$P(X = 2, Y = 3) = 0 = \frac{1}{2} \left(\frac{1}{2} - p \right)$$

$$P(X = 3, Y = 0) = 0$$

$$P(X = 3, Y = 1) = 0$$

$$P(X = 3, Y = 2) = \frac{1}{2} \left(\frac{1}{2} - p \right)$$

$$P(X = 3, Y = 3) = \frac{1}{2} \left(\frac{1}{2} - p \right)$$

Therefore

$$\begin{aligned} H(X/Y) &= -p \log_2(1) - p \log_2(1) - \frac{1}{2} \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} \right) \\ &\quad - \frac{1}{2} \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} \right) - \frac{1}{2} \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} \right) \\ &\quad - \frac{1}{2} \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} \right) \\ &= 2 \left(\frac{1}{2} - p \right) \end{aligned}$$

$$f_s = 10 \text{ kHz.}$$

$$C = \max_p [H(X) - H(X/Y)] \cdot 10^4$$

$$= \max_p \left[\left[-2p \log_2 p - 2 \left(\frac{1}{2} - p \right) \log_2 \left(\frac{1}{2} - p \right) - 2 \left(\frac{1}{2} - p \right) \right] \cdot 10^4 \right]$$

$$\frac{\partial C}{\partial p} = \left\{ -2 \log_2(p) - \frac{2p}{p} + 2 \log_2\left(\frac{1}{2} - p\right) + \frac{2\left(\frac{1}{2} - p\right)}{\left(\frac{1}{2} - p\right)} + 2 \right\} 10^4 = 0$$

which on simplification gives

$$p = 1/3$$

Therefore,

$$C = \left[-\frac{2}{3} \log_2 \frac{1}{3} - 2 \left(\frac{1}{2} - \frac{1}{3} \right) \log_2 \left(\frac{1}{2} - \frac{1}{3} \right) - 2 \left(\frac{1}{2} - \frac{1}{3} \right) \right] 10^4$$

$$C = [\log_2 3] 10^4 \text{ bits/sec.}$$

Example 2.28

If the received signal is given by

$$Y(t) = X(t) + N(t)$$

and X and Y are jointly Gaussian and the joint p.d.f. is given by

$$f_{xy}(x,y) = \frac{1}{2\pi\sigma_x\sigma_y\sqrt{1-\rho^2}} \exp \left\{ \frac{-1}{2(1-\rho^2)} \left[\left(\frac{x}{\sigma_x} \right)^2 - 2 \frac{\rho_{xy}}{\sigma_x\sigma_y} + \left(\frac{y}{\sigma_y} \right)^2 \right] \right\}$$

where

$$\rho = \frac{E(XY)}{\sigma_x\sigma_y}, \quad \sigma_x^2 = E(X^2), \quad \sigma_y^2 = E(Y^2)$$

and $X(t)$ and $N(t)$ are independent random processes. Find $I(X; Y)$ and C .

Solution

The marginal density of X is given by

$$f_x(x) = \int_{-\infty}^{\infty} f(x,y)dy = \frac{1}{(2\pi\sigma_x^2)^{\frac{1}{2}}} \exp\left(\frac{-x^2}{2\sigma_x^2}\right).$$

$$f_y(y) = \frac{1}{(2\pi\sigma_y^2)^{\frac{1}{2}}} \exp\left(\frac{-y^2}{2\sigma_y^2}\right)$$

$$\frac{f_{xy}(x,y)}{f_x(x)f_y(y)} = \frac{1}{(1-\rho_1^2)^{1/2}} \exp\left\{\frac{-\rho_1^2}{2(1-\rho_1^2)^2} \left[\left(\frac{x}{\sigma_x}\right)^2 - \frac{2xy}{\rho_1\sigma_x\sigma_y} + \left(\frac{y}{\sigma_y}\right)^2\right]\right\}$$

where $\rho_1^2 = \frac{\sigma_x^2}{\sigma_y^2}$

$$\begin{aligned} I(X, Y) &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f_{xy}(x,y) \log \frac{f_x(x)f_y(y)}{f_{xy}(x,y)} dx dy \\ &= -\frac{1}{2} \log_e(1-\rho_1^2) \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f_{xy}(x,y) dx dy \\ &\quad - \frac{\rho_1^2}{2(1-\rho_1^2)} \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f_{xy}(x,y) \left[\left(\frac{x}{\sigma_x}\right)^2 - \frac{2xy}{\rho_1\sigma_x\sigma_y} + \left(\frac{y}{\sigma_y}\right)^2 \right] dx dy \end{aligned}$$

Therefore, $I(X, Y) = -\frac{1}{2} \log_e(1-\rho_1^2)$

Since X and N are independent,

$$\sigma_y^2 = \sigma_x^2 + \sigma_N^2$$

$$\rho_1^2 = \frac{\sigma_x^2}{\sigma_y^2} = \frac{\sigma_x^2}{\sigma_x^2 + \sigma_N^2}$$

Therefore,

$$\begin{aligned} I(X, Y) &= -\frac{1}{2} \log_e \left(\frac{\sigma_N^2}{\sigma_x^2 + \sigma_N^2} \right) \\ &= \frac{1}{2} \log_e \left(\frac{\sigma_x^2 + \sigma_N^2}{\sigma_N^2} \right) \\ &= \frac{1}{2} \log_e \left(1 + \frac{\sigma_x^2}{\sigma_N^2} \right) \end{aligned}$$

Putting, $\sigma_x^2 = S$, $\sigma_N^2 = N$, noise density, we have

$$I(X, Y) = \frac{1}{2} \log_e \left(1 + \frac{S}{N} \right) \text{ nats/sample}$$

$$C = \max_{P(X)} I(X, Y) \cdot f_s$$

$$= \frac{f_s}{2} \log_e \left(1 + \frac{S}{N_0} \right)$$

$$C = B \log_e \left(1 + \frac{S}{N_0} \right) \quad \left(f_s = \frac{1}{T_s} = 2B \right)$$

where $2B$ is the bandwidth of the power spectral density of the random signal $X(t)$ and $N_0 = \eta B$ when the two-sided power spectral density of the noise is $\eta/2$ watts/Hz.

Therefore,
$$C = B \log \left(1 + \frac{S}{\eta B} \right) \text{ nats/sec.}$$

This equation is known as Shannon-Hartley theorem.

Example 2.29

A binary input-output channel is shown in fig. 2.18.

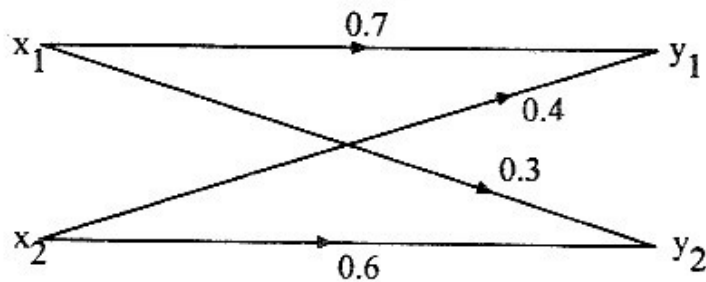


Fig. 2.18 Binary input-output channel

Given that the transition probability matrix is

$$P(Y/X) = \begin{bmatrix} 0.7 & 0.3 \\ 0.4 & 0.6 \end{bmatrix}$$

and the input probabilities are $p(x_1) = 0.5$, $p(x_2) = 0.5$. Calculate the output probability matrix $P(Y)$ and the joint probability matrix $P(X, Y)$.

Solution

$$[P(Y)] = [P(X)] [P(Y/X)]$$

$$[P(Y)] = [0.5 \quad 0.5] \begin{bmatrix} 0.7 & 0.3 \\ 0.4 & 0.6 \end{bmatrix} = [0.55 \quad 0.45]$$

If $P(X)$ is written as a diagonal matrix, then

$$[P(X, Y)] = [P(X)] [P(Y/X)]$$

$$= \begin{bmatrix} 0.5 & 0 \\ 0 & 0.5 \end{bmatrix} \begin{bmatrix} 0.7 & 0.3 \\ 0.4 & 0.6 \end{bmatrix}$$

$$[P(X, Y)] = \begin{bmatrix} 0.35 & 0.15 \\ 0.2 & 0.3 \end{bmatrix}$$

Review Questions

- 2.1 What do you mean by memoryless channel?
- 2.2 Define a discrete channel.
- 2.3 When is a discrete channel said to be “memoryless”?
- 2.4 Define Mutual information / Transinformation
- 2.5 Distinguish between noisy reception and perfect reception.
- 2.6 Name the different types of channels.
- 2.7 Define channel matrix D.
- 2.8 What is a lossless channel?
- 2.9 Define a deterministic channel.
- 2.10 What is a noiseless channel?
- 2.11 When is a channel said to be useless?
- 2.12 Define a symmetric channel.
- 2.13 What do you understand by BSC and BEC.
- 2.14 Define channel capacity.
- 2.15 What is the channel capacity of (i) lossless channel (ii) Deterministic channel (iii) Noiseless channel (iv) Symmetric channel.

- 2.16 How do you find the channel capacity of unsymmetric channels?
- 2.17 What is the channel capacity of (i) BEC (ii) BSC.
- 2.18 Write down Fano's Inequality.
- 2.19 State the Shannon's Fundamental theorem / Noisy coding Theorem.
- 2.20 What do you mean by a gaussian channel?
- 2.21 How do you find the capacity of a gaussian channel?
- 2.22 State Shannon-Hartely's Law.
- 2.23. Define information rate of the source.
- 2.24 Given a channel matrix, how do you find the channel capacity?
- 2.25 State the applications of coding Techniques.
- 2.26 State the advantages of coding techniques.
- 2.27 Distinguish between the different types of channels.
- 2.28 How do you obtain the channel capacity of a symmetric channel.
- 2.29 Define a stochastic matrix.
- 2.30 Distinguish between discrete and continuous type channels.